

# Choice offering imperatives in inquisitive and truth-maker semantics

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Imperatives: worlds and beyond  
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# Part I

## Setting the stage

- ▶ How can we understand the meaning of an imperative sentence?
- ▶ In the case of a declarative sentence, there is a traditional answer: truth-conditions.
- ▶ Semantics comes in the form of a truth relation  $w \models \varphi$  defined between “possible worlds” and sentences.
- ▶ Let  $|\varphi| := \{w \mid w \models \varphi\}$
- ▶ Then, a declarative can be seen as partially describing a possible world by locating it inside the set  $|\varphi|$ .

Natural hypothesis: imperative meaning = **compliance conditions**.

Let us try to implement this idea in a basic setting: propositional logic.

**Action formulas:**  $\varphi ::= a \mid \neg\varphi \mid \varphi \wedge \varphi \mid \varphi \vee \varphi$

- ▶  $a \rightsquigarrow$  calling Alice
- ▶  $\neg a \rightsquigarrow$  not calling Alice
- ▶  $a \wedge b \rightsquigarrow$  calling Alice and calling Bob
- ▶  $a \vee b \rightsquigarrow$  calling Alice or calling Bob

**Imperatives** are formulas of the form  $!\varphi$ , for  $\varphi$  an action formula.

- ▶ We take compliance conditions to be assessed relative to a **conduct**.
- ▶ Intuitively, a conduct encodes what actions the relevant agent has performed over the relevant stretch of time.
- ▶ Formally, a conduct is just a valuation function  $c : \mathcal{A} \rightarrow \{0, 1\}$ .  
 $c(a) = 1$  if the agent performed some  $a$ -actions,  $c(a) = 0$  otherwise.

## Execution conditions

- ▶  $c \models a \iff c(a) = 1$
- ▶  $c \models \neg\varphi \iff c \not\models \varphi$
- ▶  $c \models \varphi \wedge \psi \iff c \models \varphi \text{ and } c \models \psi$
- ▶  $c \models \varphi \vee \psi \iff c \models \varphi \text{ or } c \models \psi$

## Compliance conditions

- ▶  $c \models !\varphi \iff c \models \varphi$

Compliance set:  $|\varphi| := \{c \mid c \models \varphi\}$

**Intuition:** like a declarative describes a possible world, an imperative describes a conduct by locating it within  $|\varphi|$ .

## Imperative entailment

$$!\varphi \models !\psi \iff \forall c : c \models !\varphi \Rightarrow c \models !\psi$$

Prediction: the logic of imperatives is classical

$$!\varphi \models !\psi \iff \varphi \text{ entails } \psi \text{ in classical logic.}$$

If this were right, the difference between the two kinds of sentences would lie only in what they describe: states of affairs vs. ways to act.

However, this prediction seems contradicted by basic examples.

## Problem I: Ross's paradox

As Ross ('41) observed, the entailment  $!a \models !(a \vee b)$  fails for imperatives. When told (1-a), one is not entitled to infer (1-b).

- (1) a. Post this letter.
- b. Post this letter or burn it.

Yet, complying with (1-a) implies complying with (1-b)!

The contrapositive entailment  $!\neg(a \vee b) \models !\neg a$  does seem valid. From (2-a), one is entitled to infer (2-b).

- (2) a. Do not post or burn this letter.
- b. Do not post this letter.

## Problem II: Veltman's puzzle

Imagine that a patient consults two doctors about the same problem and obtains the following two prescriptions.

- (3)     a.    Drink milk or apple juice.  
          b.    Don't drink milk.

- ▶ Intuitively, the doctors are disagreeing, and cannot both be trusted.
- ▶ In terms of compliance, though, the prescriptions are consistent. Our patient should simply draw the conclusion: drink apple juice.
- ▶ We would like our account to explain why (3-a) and (3-b) are perceived as conflicting, though it is possible to comply with both.

## Part II

### The basic account

Why does (4-a) not entail (4-b)?

- (4)    a.    Post this letter.  
      b.    Post this letter or burn it.

### Common intuition:

if I tell you to post the letter, I'm not giving you the option to burn it.

- ▶ We take this intuition seriously as the starting point of our proposal.
- ▶ To make sense of this as an explanation of the non-entailment, we have to assume that (4-b) does grant an option to burn the letter.
- ▶ The meaning of (4-b) has a permission component that is not present in (and in fact is contradicted by) the content of (4-a).
- ▶ This explains why (4-a) is not stronger than (4-b).

The same intuition explains Veltman's puzzle.

- (5)    a.    Drink milk or apple juice.  
       b.    Don't drink milk.

The doctors are disagreeing because:

- ▶ (5-a) implies that milk is an option;
- ▶ (5-b) implies that milk is not an option.



## The meaning of an imperative

An imperative  $!\varphi$  conveys two things about the option set  $s$ :

- ▶ all options in  $s$  execute  $\varphi$  in some way;
- ▶ for any way of executing  $\varphi$  there is a corresponding option.

To formalize this, need a semantics for action formulas that tells us what the ways of executing a given action formula.

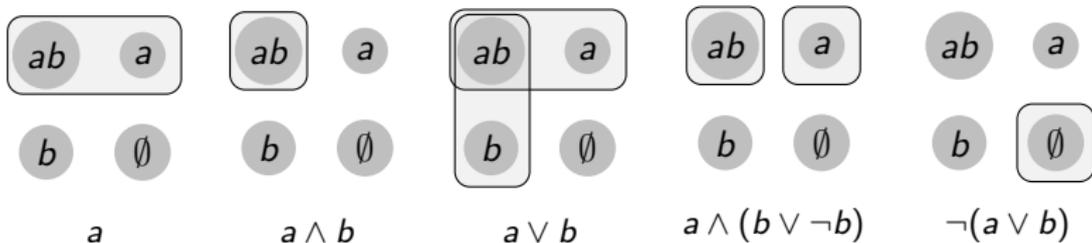
We move from compliance-conditional semantics to **inquisitive semantics**

## Inquisitive semantics for action formulas

The meaning of an action formula is given by support at an option set. We can read  $s \models \varphi$  as:  $\varphi$  is executed uniformly in all the conducts in  $s$ .

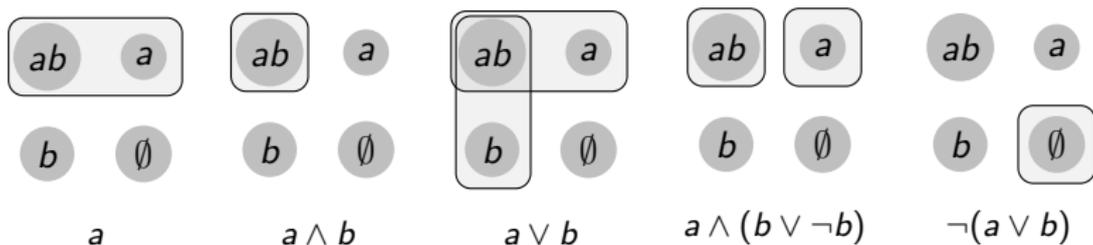
- ▶  $s \models a \iff c(a) = 1$  for all  $c \in s$
- ▶  $s \models \varphi \vee \psi \iff s \models \varphi$  or  $s \models \psi$
- ▶  $s \models \varphi \wedge \psi \iff s \models \varphi$  and  $s \models \psi$
- ▶  $s \models \neg\varphi \iff s \cap t = \emptyset$  for all  $t \models \varphi$

Let  $\text{Alt}(\varphi)$  be the set of maximal option sets that support  $\varphi$ .



Definition:

we call  $!\varphi$  a **basic imperative** if  $\text{Alt}(\varphi)$  is a singleton,  
and a **choice imperative** if it contains multiple elements.



Fact 1:  $c \models \varphi \iff \{c\} \models \varphi$

Fact 2:  $|\varphi| = \bigcup \text{Alt}(\varphi)$

## Semantics of imperatives:

Writing  $s \not\sim t$  for  $s \cap t \neq \emptyset$ , we let:

$$s \models !\varphi \iff \begin{array}{l} (i) \quad s \subseteq |\varphi| \\ (ii) \quad s \not\sim \alpha \text{ for all } \alpha \in \text{Alt}(\varphi) \end{array}$$

## Semantics of imperatives, restated:

$$s \models !\varphi \iff \begin{array}{l} (i) \quad \forall c \in s \quad \exists \alpha \in \text{Alt}(\varphi) : c \in \alpha \\ (ii) \quad \forall \alpha \in \text{Alt}(\varphi) \quad \exists c \in s : c \in \alpha \end{array}$$

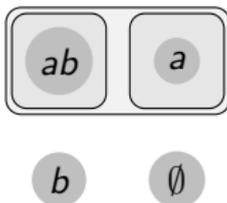
## Imperative entailment and equivalence:

- ▶  $!\varphi \models !\psi \iff \forall s : s \models !\varphi \Rightarrow s \models !\psi$
- ▶  $!\varphi \equiv !\psi \iff !\varphi \models !\psi \text{ and } !\psi \models !\varphi$

## Predictions I: basic imperatives

(6) Call Alice.  $\rightsquigarrow$  !a

- ▶ Atomic imperatives are basic:  $\text{Alt}(a) = \{|a|\}$
- ▶  $s \models !a \iff s \subseteq |a| \text{ and } s \not\subseteq |a| \iff s \subseteq |a| \text{ and } s \neq \emptyset$



**Prediction:** (6) conveys that all options are conducts in which A is called.

## Predictions I: basic imperatives

What we just saw for atomic imperatives extends to all basic imperatives.

- ▶ Let  $\varphi$  be a basic imperative, i.e.,  $\text{Alt}(\varphi) = \{|\varphi|\}$ .
- ▶  $s \models !\varphi \iff s \subseteq |\varphi|$  and  $s \neq \emptyset$

**Prediction:**  $!\varphi$  conveys that all options are conducts that execute  $\varphi$ .  
The meaning of  $!\varphi$  is fully determined by the execution conditions of  $\varphi$ .

**Prediction:**  $!\varphi \models !\psi \iff |\varphi| \subseteq |\psi| \iff \varphi \models_{\text{CPL}} \psi$   
The logic of basic imperatives is classical.

## Predictions II: disjunctive imperatives

(7) Call Alice or Bob.  $\rightsquigarrow$   $!(a \vee b)$

- ▶  $\text{Alt}(a \vee b) = \{|a|, |b|\}$ .
- ▶  $s \models !(a \vee b) \iff s \subseteq |a \vee b|$  and  $s \not\subseteq |a|$  and  $s \not\subseteq |b|$

**Prediction:** (7) conveys that

1. in all the available options, either A or B is called;
2. there is an option in which A is called, and one in which B is called.

## Predictions II: Ross's observation accounted for

- ▶ Prediction:  $\neg a \not\models \neg(a \vee b)$

If  $c(a) = 1$  and  $c(b) = 0$ , then  $\{c\} \models \neg a$  but  $\{c\} \not\models \neg(a \vee b)$   
Hence, the inference from (8-a) to (8-b) is predicted as invalid.

- (8) a. Call Alice.  
b. Call Alice or Bob.

- ▶ In the [Ross scenario](#), moreover, posting and burning are exclusive. Taking this into account,  $\neg p$  and  $\neg(p \vee b)$  are in fact inconsistent.



- ▶ Thus, far from following from the imperative *Post this letter*, the imperative *Post this letter or burn it* contradicts it.

## Predictions III: conjunctive imperatives

- (9)    a.    Call Alice and Bob.  $!(a \wedge b)$   
      b.    Call Alice, and call either Bob or Charlie.  $!(a \wedge (b \vee c))$

- ▶ (9-a) is a basic imperative:  $\text{Alt}(a \wedge b) = \{|a \wedge b|\}$
- ▶ **Prediction:** (9-a) conveys an obligation to call both.
- ▶ Notice that (9-a) is correctly predicted to entail both  $!a$  and  $!b$ .
  
- ▶ (9-b) is a choice imperative:  $\text{Alt}(a \wedge (b \vee c)) = \{|a \wedge b|, |a \wedge c|\}$ .
- ▶  $s \models !(a \wedge (b \vee c)) \iff s \subseteq |a| \text{ and } s \subseteq |b \vee c| \text{ and } s \not\subseteq |b| \text{ and } s \not\subseteq |c|$
- ▶ **Prediction** (9-b) conveys that:
  - ▶ every option is a conduct in which A is called;
  - ▶ every option is a conduct in which at least one of B and C is called;
  - ▶ there is an option in which B is called, and one in which C is called.

## Predictions IV: negative imperatives

- (10)    a.    Don't post this letter.  $!\neg p$   
      b.    Don't post or burn this letter.  $!\neg(p \vee b)$

- ▶ Negative imperatives are always basic:  $\text{Alt}(\neg\varphi) = \{\overline{|\varphi|}\}$
- ▶ **Prediction:** a negative imperative  $!\neg\varphi$  conveys that every option in  $s$  is a conduct in which  $\varphi$  is not executed.
- ▶ (10-b) is correctly predicted to entail (10-a):  $!\neg(p \vee b) \models !\neg p$ .  
This because both imperatives are basic and  $\overline{|p \vee q|} \subseteq \overline{|p|}$ .
- ▶ Thus, our observations about Ross's example are accounted for.

## Predictions IV: Veltman's puzzle accounted for

- (11) a. Drink milk or apple juice.  
b. Don't drink milk.

$$!(m \vee a) \\ !\neg m$$

- ▶ These two imperatives are indeed inconsistent:
  - ▶ if  $s \models !(m \vee a)$ , then  $s \not\subseteq |m|$
  - ▶ but if  $s \models !\neg m$ , then  $s \subseteq \overline{|m|}$
  - ▶ no  $s$  can satisfy both conditions
- ▶ This explains the feeling that the prescriptions are conflicting, in spite of the fact that it is possible to comply with both.

## Taking stocks

- ▶ We have proposed an account in which an imperative describes a set of options for action.
- ▶ In particular, we have assumed that an imperative  $!\varphi$ :
  - ▶ provides an obligation to comply with  $\varphi$ ;
  - ▶ offers a choice between alternative ways to comply.
- ▶ We have shown that this offers a natural explanation of two interesting puzzles concerning the logic of imperatives:
  - ▶ Ross's paradox
  - ▶ Veltman's puzzle

# Part III

## Exhaustivity

## Hurford's constraint

- ▶ Hurford (1974) observed that disjunctions in which one disjunct entails the other are generally infelicitous:

- (12)
- a. #John is an American or a Californian.
  - b. #The number  $x$  is greater than 6 or different from 6.
  - c. #Alice came, or Alice came and Bob didn't.

- ▶ This is commonly referred to as Hurford's constraint.

- ▶ Hurford's constraint is also at play in imperatives:

- (13)
- a. #Hire an American or a Californian.
  - b. #Take a number greater than 6 or different from 6.
  - c. #Call Alice, or call Alice and don't call Bob.

## Why Hurford's constraint?

(14) #Hire an American or a Californian.

- ▶ Hurford's constraint has been explained in terms of a more general ban against LFs which contain redundant operators.

(Katzir and Singh 2013, Simons 2001, Meyer 2013)

- ▶ Our account allows us to extend the explanation to imperatives:
  - ▶ if  $\psi \models \varphi$ , then in inquisitive semantics  $\varphi \vee \psi \equiv \varphi$ ;
  - ▶ as a consequence, the Hurford disjunctive imperatives in (14) contain structural redundancy, and are predicted to be deviant.

- ▶ There are cases that seem to violate Hurford's constraint:

(Gazdar 1979)

- (15)
- a. Alice called Bob, Charlie, or both.
  - b. Alice took two or three cards.

- ▶ The situation is analogous for imperatives:

- (16)
- a. Call Bob, Charlie, or both.
  - b. Take two or three cards.

- ▶ How to account for these cases?

- ▶ To deal with these cases, in AC13 we used a semantics in which  $\text{Alt}(p \vee q) = \{|p|, |q|\}$  even when  $|q| \subseteq_c |p|$ .
- ▶ *Mutatis mutandis*, the same holds for Fine's truth-maker account.
- ▶ However, this solution creates two other problems.
- ▶ First, in these systems  $\varphi \vee \psi \not\equiv \varphi$  even when  $\psi \models \varphi$ . So, we lose our explanation for the oddness of (17-a,b).

- (17)
- #Hire an American or a Californian.
  - #Call Alice, or call Alice and not Bob.

- ▶ We would like a theory that explains the difference between “good” and “bad” Hurford disjunctions.

- ▶ Moreover, consider the following sentences:

(18)    a.    Call Alice, or call Alice, Bob, and Charlie.  
          b.    Take two cards or four cards.

- ▶ Intuitively, calling just Alice and Bob is a contravention to (18-a), and taking exactly three cards is a contravention to (18-b).
- ▶ But both our AC13 account and Fine's account wrongly predict such conducts to be (entirely) compliant.
- ▶ To solve both problems at once we adopt the explanation proposed by Chierchia, Fox and Spector (2009) for declaratives.

► **Observation:**

the felicitous disjunctions are those in which the weak disjunct admits an exhaustive interpretation that breaks the entailment.

- (19) a. Call Alice  $\rightsquigarrow$  not Bob, Charlie, ...  
b. Take two cards  $\rightsquigarrow$  not three

- (20) a. Hire an American  $\not\rightsquigarrow$  not Californian  
b. Call Alice  $\not\rightsquigarrow$  not (Alice and not Bob)

► **Proposal:**

this strengthening can sometimes occur in embedded positions.

- Concretely, this happens by means an optional operator *exh*, which amounts to a silent *only*.

## Good Hurford disjunctions

(21) Call Alice, Bob, or both.

- ▶ Hurford's constraint rules out the LF  $a \vee b \vee (a \wedge b)$
- ▶ however, we also have the LF  $\text{exh}(a) \vee \text{exh}(b) \vee \text{exh}(a \wedge b)$
- ▶ under natural assumptions this is  $(a \wedge \neg b) \vee (b \wedge \neg a) \vee (a \wedge b)$
- ▶ this LF no longer violates Hurford's constraint;
- ▶ since it has a suitable LF, (21) is felicitous.

## Bad Hurford disjunctions

(22) #Call Alice or call Alice but not Bob.

- ▶ HC rules out the LF  $a \vee (a \wedge \neg b)$
- ▶ we also have the LF  $\text{exh}(a) \vee \text{exh}(a \wedge \neg b)$
- ▶ under natural assumptions, this amounts to  $(a \wedge \neg b) \vee (a \wedge \neg b)$
- ▶ however, this LF still violates Hurford's constraint;
- ▶ since (22) has no suitable LF, (22) is infelicitous.

# Exhaustification: definitions

The *exh* operator

$s \models \text{exh}(\varphi) \iff s \subseteq \text{exh}(\alpha, |RA(\varphi)|)$  for some  $\alpha \in \text{Alt}(\varphi)$ .

Where:

- ▶  $|RA(\varphi)| = \{|\psi| \mid \psi \in RA(\varphi)\}$
- ▶  $\text{exh}(\pi, \Pi) = \pi - \bigcup \{\pi' \in \Pi \mid \pi \not\subseteq \pi'\}$
- ▶ The set  $RA(\varphi)$  is defined recursively as follows:

$$\begin{aligned}RA(a) &= \{a\} \cup C_a \\RA(\varphi \vee \psi) &= RA(\varphi) \cup RA(\psi) \\RA(\varphi \wedge \psi) &= RA(\varphi) \cup RA(\psi) \\RA(\neg\varphi) &= \{\neg\psi \mid \psi \in RA(\varphi)\} \\RA(\text{exh}(\varphi)) &= \{\text{exh}(\psi) \mid \psi \in RA(\varphi)\}\end{aligned}$$

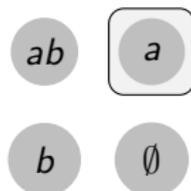
where  $C_a$  is a set of contextually relevant alternatives to  $a$ .

# Illustration: atomic imperatives

(23) Call Alice.

►  $\text{exh}(a) \equiv (a \wedge \neg b)$

$C = \{a, b\}$



► Prediction: ambiguity

(24) a.  $!a$   
b.  $!\text{exh}(a)$

$\not\Rightarrow !\neg b$   
 $\Rightarrow !\neg b$

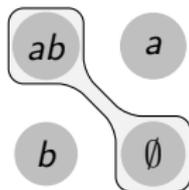


## A problem: exh under negation

(27) Don't call Alf or Bea.

►  $\neg \text{exh}(a \vee b) \equiv \neg((a \wedge \neg b) \vee (b \wedge \neg a))$

$$C = \{a, b\}$$



- (28) a.  $!\neg(a \vee b)$   
b.  $!\neg \text{exh}(a \vee b)$

[call neither]  
[don't call exactly one]

► Wrong Prediction!

# Constraints on the application of exh

- ▶ The application of exh is optional, but **constrained**:

(29) exh can only apply when it creates a stronger meaning or is needed to avoid violations of pragmatic constraints.

- ▶ Consequence: only (30) is ambiguous

(30) Call Alf or Bea.

- a.  $!(a \vee b)$
- b.  $!\text{exh}(a \vee b)$

(31) Don't call Alf or Bea.

- a.  $!\neg(a \vee b)$
- b.  $\#\neg\text{exh}(a \vee b) \rightsquigarrow$  ruled out by (29)

# Distinguishing good and bad Hurford's cases

- ▶ Local exh can obviate Hurford's violations:

(32) Call Alice, or Bob, or both.

a.  $!(a \vee b \vee (a \wedge b))$  [violation]

b.  $!(\text{exh}(a) \vee \text{exh}(b) \vee \text{exh}(a \wedge b)) \equiv$   
 $!((a \wedge \neg b) \vee (b \wedge \neg a) \vee (a \wedge b))$  [ok]

- ▶ But this does not always succeed:

(33) #Call Alice or call Alice but not Bob.

a.  $a \vee (a \wedge \neg b)$  [violation]

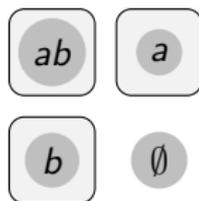
b.  $\text{exh}(a) \vee \text{exh}(a \wedge \neg b) \equiv$   
 $(a \wedge \neg b) \vee (a \wedge \neg b)$  [violation]

# Interpreting good Hurford imperatives

(34) Call Alice, or Bob, or both.

$$C = \{a, b\}$$

- ▶ The only acceptable logical form for (34) is:



$$\text{exh}(a) \vee \text{exh}(b) \vee \text{exh}(a \wedge b)$$

- ▶ **Prediction:** (34) conveys a strong free choice:
  - ▶ Calling only A is an option.
  - ▶ Calling only B is an option.
  - ▶ Calling both is an option.

# Predicting compliance gaps

(35) Call Alice, or call Alice, Bob, and Charlie.

$C = \{a, b, c\}$

- ▶ The only acceptable logical form for the action formula in (35) is:



$\text{exh}(a) \vee \text{exh}(a \wedge b \wedge c)$

- ▶ **Prediction:**
  - ▶ Addressee must either call only A, or call A, B, and C.
  - ▶ There is an option to call only A.
  - ▶ There is an option to call A, B, and C.
- ▶ In particular, we predict (35)  $\models$  *not only Alice and Bob*.

## Wrapping up

- ▶ We complement our base theory with an optional, but constrained, exhaustification operator.
- ▶ This gives us a principled explanation of the difference between:  
(36) a. Take one or two cards.  
b. #Take a card or an ace.
- ▶ Predicts subtle facts on the interpretation of Hurford imperatives:  
(37) Take one or three cards.  
↪ not exactly two

# Part IV

## Discussion

# Imperatives in inquisitive and truth-maker semantics

|                     | CA16     | AC13     | Fine15   |
|---------------------|----------|----------|----------|
| free-choice         | yes      | yes      | yes      |
| hyperintensional    | yes      | yes      | yes      |
| semantic primitives | conducts | conducts | actions  |
| kind of semantics   | inexact  | exact    | exact    |
| exhaustivity        | local    | global   | none     |
| entailment          | standard | standard | analytic |
| FC atoms            | no       | no       | yes      |
| FC negations        | no       | no       | yes      |
| de Morgan           | invalid  | invalid  | valid    |

# Can atomic imperatives be choice-offering?

- ▶ In our semantics, we chose to make atomic imperatives basic.
- ▶ In Fine15, atomic imperatives may be choice offering.
- ▶ If we take atoms to stand for logically simple imperative sentence, this is an interesting empirical issue.
- ▶ This issue teases apart two hypotheses on the sources of free-choice:
  - ▶ **ontic**: choice stems from the existence of multiple compliant actions;
  - ▶ **grammatical**: choice is introduced by certain linguistic devices.
- ▶ To clarify this, we can try to simulate a disjunctive imperative by an atomic one which is intuitively as close as possible to the disjunction.

# Can atomic imperatives be choice-offering?

- (38) a. Come visit us in the weekend.  
b. Come visit us on Saturday or Sunday.

In spite of having (apparently) the same compliant actions, we think that only (38-b), but not (38-a), is choice-offering.

## Argument 1

- (39) MOTHER Come and visit us next Saturday or Sunday.  
FATHER Don't come on Sunday.  
SON But mom told me I *could* come on Sunday.
- (40) MOTHER Come and visit us next weekend.  
FATHER Don't come on Sunday.  
SON ??But mom told me I *could* come on Sunday.

## Argument 2

(41) Come visit us next Sunday!

When told (41), (42-a) seems a truthful report, but (42-b) is dubious.

- (42) a. They invited me to visit them next weekend.  
b. ?They invited me to visit them next Saturday or Sunday.

Assuming that a report of  $\psi$  is truthful only if  $\psi$  follows from the original imperative, this indicates that (43-b) follows from (43-a).

- (43) a. Come visit us next Saturday.  
b. Come visit us next weekend.

This is predicted is (43-b) is basic, but not if it is a choice imperative.

This suggests that, even when “disjunctive” in nature, atomic imperatives cannot be choice-offering, and that choice has a grammatical source.

# Can negative imperatives be choice-offering?

- ▶ In our semantics, negative imperatives are always basic. There is only one way to realize  $\neg\varphi$ : abstain from realizing  $\varphi$ .
- ▶ In Fine15, instead, negative imperatives may be choice offering.
- ▶ In particular, a negation of a conjunction like (44-a) is choice-offering, and equivalent to (44-b).

- (44)    a.    Don't drink and drive.  
          b.    Don't drink or don't drive.

- ▶ So, is (44-a) choice-offering?

- ▶ We think this is not the case:

(45) a. Don't drink.  
b. So, don't drink and drive.

(46) a. Don't drink.  
b. ??So, don't drink or don't drive.

- ▶ This supports our prediction that only (46-b) is choice-offering, while (45-b) is basic and has classical logical properties.
- ▶ This supports the view that negative imperatives are always basic.

## Conducts vs. actions

- ▶ We believe that negative imperatives also expose an advantage of a semantics based on conducts as opposed to actions.
- ▶ It is not clear what action is to count as exactly complying with (47).

(47) Don't talk to Alice.

- ▶ It seems that one complies with (47) not by performing some action, but rather by not performing any action of a certain kind.
- ▶ In other words, compliance with (47) is not a property of any one action in a conduct, but rather a property of the conduct as a whole.
- ▶ On our approach, the meaning (47) is simple: (47) conveys an obligation to keep one's conduct among those in which there is no action of talking to Alice.
- ▶ We believe that this is important to assess the theory's predictions.

# Idempotency

- ▶ Both AC13 and Fine15 make conjunction not idempotent:

- (48)
- a. Take a card.  $\neq$
  - b. Take a card and take a card.

(48-b) conveys an obligation to take at least a card, and a permission to take any card and to take **any two cards**.

- ▶ In CA16, conjunction is idempotent, and so the oddness of (48-b) is explained by the ban on structural redundancy.
- ▶ In AC13 and Fine15, the conjunction in (48-b) is not redundant, and one is left wondering why (48-b) is not usable instead of (49).

- (49) Take a card or take two cards.

# Mereological entailment

- ▶ Fine15 accounts for imperative entailment by a mereological notion.
- ▶ However, this notion is vulnerable to a version of Ross's paradox: (50-a) is predicted to entail (50-b).

- (50) a. Take one or three cards.  
b. Take one, two, or three cards.



- ▶ In fact, the opposite entailment holds as well, which shows that in this approach, logical equivalence does not guarantee synonymy.

In our account, the only LF for these sentences involves exh:

- (51) a. Take one or three cards.  $\rightsquigarrow \text{exh}(1) \vee \text{exh}(3)$   
b. Take one, two, or three cards.  $\rightsquigarrow \text{exh}(1) \vee \text{exh}(2) \vee \text{exh}(3)$



Neither entailment is valid:

- ▶ (a)  $\not\models$  (b) since (a) does not give permission to take exactly two;
- ▶ (b)  $\not\models$  (a) since (b) does not forbid taking exactly two.

## Conclusions

- ▶ We have proposed that imperatives describe sets of options.
- ▶ They do not just specify restrictions on the available conduct, but can also convey the existence of certain options.
- ▶ Key features of our theory:
  - ▶ an inquisitive semantics of action formulas;
  - ▶ an alternative-sensitive imperative operator !;
  - ▶ optional exhaustification operator.
- ▶ The theory provides natural explanations for a number of puzzles concerning the logic of imperatives.

## Further work

- ▶ Conditional imperatives.
- ▶ Free choice stemming from indefinites (take any card).
- ▶ Connection with deontics.

