

# Diego's Theorem

Dick de Jongh/ ILLC, Universiteit van Amsterdam

February 28, 2017

**Theorem 1** (Diego,1961(!)). The fragment  $[\rightarrow]$  of formulas of IPC with only  $\rightarrow$  is locally finite, i.e. the number of equivalence classes of formulas in a finite number of variables is finite.

The proof takes a number of steps. The proof is sketched for formulas in  $[\rightarrow, \wedge, \perp]$ , which is not more difficult. An  $n$ -model is a Kripke model for the propositional variables  $p_0, \dots, p_{n-1}$ .

**Definition 2.** A world  $w$  in a Kripke model  $\mathfrak{M}$  is an *intersection node* if  $w$  is not terminal, and for each proposition letter  $p$ ,  $w \in V(p)$  iff  $w' \in V(p)$  for each  $w'$  with  $wR^+w'$ .

We write  $uR^+v$  for  $uRv$  and  $u \neq v$ .

This means that, if  $p$  fails at an intersection node it fails higher up in the model.

**Proposition 3.** If  $w$  is an intersection node, then, for all formulas  $\varphi$  in  $[\rightarrow, \wedge, \perp]$ ,  $w \Vdash \varphi$  iff  $w' \Vdash \varphi$  for each  $w'$  with  $wR^+w'$ .

*Proof.* By induction on the length of  $\varphi$ . We just do the cases of  $\wedge$  and  $\rightarrow$ .

$w \Vdash \psi \wedge \chi$  iff  $w \Vdash \psi$  and  $w \Vdash \chi$  iff (by IH)  $w' \Vdash \psi$  for each  $w'$  with  $wR^+w'$  and  $w' \Vdash \chi$  for each  $w'$  with  $wR^+w'$  iff  $w' \Vdash \psi \wedge \chi$  for each  $w'$  with  $wR^+w'$ .

$w \not\Vdash \psi \rightarrow \chi$  iff ( $w \Vdash \psi$  and  $w \not\Vdash \chi$ ) or  $\exists w'(wR^+w'$  and  $w' \Vdash \psi$  and  $w' \not\Vdash \chi$ ) iff (by IH) ( $w \Vdash \psi$  and  $\exists w'(wR^+w'$  and  $w' \not\Vdash \chi$ ) or  $\exists w'(wR^+w'$  and  $w' \Vdash \psi$  and  $w' \not\Vdash \chi$ ). In both cases this implies  $\exists w'(wR^+w'$  and  $w' \not\Vdash \psi \rightarrow \chi$ ), and the other direction is trivial. □

**Proposition 4.** If  $\mathfrak{M}$  is a finite rooted Kripke model and  $\mathfrak{M}'$  arises from  $\mathfrak{M}$  by deleting all intersection nodes but the root, then, for all formulas  $\varphi$  in  $[\rightarrow, \wedge, \perp]$ , and all  $w$  in  $\mathfrak{M}'$ ,  $w \Vdash \varphi$  iff  $\mathfrak{M}, w \Vdash \varphi$ .

*Proof.* Induction on the length of  $\varphi$ . The case of  $\varphi = \psi \rightarrow \chi$  runs as follows: Take  $w$  in  $\mathfrak{M}'$ .  $\mathfrak{M}, w \not\Vdash \varphi$  iff, (for some non-intersection node  $w'$  with  $wRw'$ ,  $\mathfrak{M}, w' \Vdash \psi$  and  $\mathfrak{M}, w' \not\Vdash \chi$ ) or (for some intersection node  $w'$  with  $wRw'$ ,  $\mathfrak{M}, w' \Vdash \psi$  and  $\mathfrak{M}, w' \not\Vdash \chi$ ). Now, the latter is, by repeatedly applying Proposition ?? noting that the model is finite and intersection nodes on non-terminal, to: for some intersection node  $w'$  with  $wRw'$ ,  $\mathfrak{M}, w' \Vdash \psi$  and  $\mathfrak{M}, w' \not\Vdash \chi$  iff, by induction hypothesis,  $\mathfrak{M}', w \not\Vdash \varphi$ . □

Note that an  $n$ -model without intersection nodes the number of atoms forced in a chain running from top to bottom decreases at each step, so that chains cannot be longer than  $n$ , i.e. the depth of an  $n$ -model without intersection nodes is bounded by  $n$ . We have now proved that, if two  $n$ -formulas  $\varphi, \psi \in [\rightarrow, \wedge, \perp]$  are not IPC-equivalent and hence can be distinguished on a finite model, then they can be distinguished on a model of depth  $n$  or smaller.

We now sketch a technique which will be made more precise when we discuss universal models to obtain the fact there is a finite set of models that distinguishes all formulas in  $[\rightarrow, \wedge, \perp]$ , and hence this class of formulas contains only a finite number of equivalence classes. We need think only of the case that a model distinguishes two formulas in its root. Also, we need consider only models that are trees.

We call a function from  $(p_0, \dots, p_{n-1})$  to  $\{0, 1\}$  a *color*. There are  $2^n$  such colors, so only  $2^n$  non-isomorphic rooted  $n$ -models of depth 1. Suppose we have a class  $X$  of  $n$ -models of depth  $m$

that suffices to distinguish all pairs of formulas that can be distinguished on a model of depth  $m$ . Then a class of models of depth  $m + 1$  can be obtained with the same property by putting a root with a fitting color below a subset of  $X$ .